**Optimizing A\* with lookahead and graph pruning**

We are often faced with the problem of finding an optimal path between two points on a uniform grid consisting of wall tiles and empty tiles where only orthogonal movement is permitted. A special implementation of A\* can be created to take advantage of the grid-structure of our graph. Furthermore, because we are on a uniform grid, we can perform two different optimizations that can speed up the search. The first optimization we consider is inspired by jump point search (JPS) [1]. JPS is designed for diagonally-connected grids, but we can use a similar method to optimize our pathfinding on an orthogonally-connected grid, we call this optimization lookahead. The second optimization we can make is graph pruning using cellular automata [2]. We compare performance between A\*, lookahead, pruned A\*, and pruned lookahead.

**Lookahead**

Our typical A\* algorithm involves storing a heap of unsearched nodes, and repeatedly popping the node with the lowest f value, which is the sum of the distance travelled (g) and the heuristic (h). our heuristic is the Manhattan distance to the destination, this is admissible on a uniform grid where all costs are 1. When we pop a node of cost f, we check its neighbours and add each unsearched neighbour to the heap. In lookahead, we can make the following optimizations:

1. We can skip checking the parent node, since we have already visited it.
2. If a node diagonally “behind” us is empty and unsearched, we can skip the nodes on that “side”, since the parent node has an equally good path to these nodes. The reason the diagonal node must be unsearched, is because it may
3. If any node adjacent to the current node has a smaller heuristic, we can choose one of these nodes to immediately check after the current node. This works because we know that fcurrent was a minimum value of the heap, and for any child fchild = gchild + hchild = (gcurrent + 1) + (h­current ± 1) = fcurrent + 1 ± 1, so fchild >= fcurrent, so if we find a node such that fchild = fcurrent, we know that it must be a minimum value of the heap, so we can safely choose it as the next node to search without adding it to the heap.

Figure 1 – lookahead optimizations 1, 2 and 3 illustrated on a 2d grid, these are trivially generalizable to n dimensions.

Optimization 2 has a problem: it cannot always be applied. Consider

**Cellular automata pruning**

Given the dimension of the grid (d) we can prune the graph with two simple rules:

1. If a cell is a dead end (2d - 1 orthogonally adjacent walls), we can safely prune it.
2. If a cell is a corner (d orthogonally adjacent walls that are themselves diagonally adjacent), and part of a 2d grid of empty cells, we can safely prune it.

Since any tile outside the grid is impassable, we can count these as walls for the purposes of pruning. Additionally, we can count pruned tiles as walls for the purposes of pruning, which allows further pruning.

Figure 2 – rules for pruning illustrated on a 2d grid

The pseudocode for our pruning algorithm looks like this:

Prune(graph, coord):

If cell at coord is already a wall or pruned, or it is the start or end cell, return

Get the diagonally adjacent neighbours of coord

If should\_prune(neighbours):

Mark cell at coord as pruned

Call Prune on all diagonally adjacent neighbours of coord

Then we simply call Prune on every cell.

Since each neighbour contains one bit of information (wall/pruned or not), we can optimize should\_prune by packing the neighbours into an index and using a lookup table.

Figure 3 – bit-indexes of each adjacent cell

We use some rust code to generate the lookup table for a 2d grid:

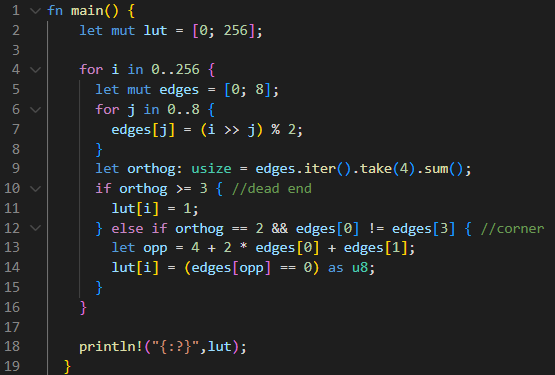


Figure 4 lookup table calculation

In our analysis, we know our start and end points ahead of time, so we can avoid pruning them. In another scenario, we might not know how to avoid pruning away our start and end points. If we prune away the start or end point, we can solve it in two ways, either we can use floodfill to unprune all pruned nodes around the point, or we can permit the pathfinder to move from a pruned node to another pruned node.

**Complexity Analysis**

We assume the heuristic calculation is O(1), and so not a factor in time complexity calculations. In the worst case on an arbitrary graph, A\* is equivalent to dijkstra’s algorithm and must search O(bd)[3] nodes, where b is the branching factor of the graph and d is the distance from the starting node to the goal. In our application, the graph is known to be a uniform grid, which reduces the worst case performance to O(dn) nodes, where n is the dimensionality of our grid. On an arbitrary graph, the average case time complexity is O(b\*d)[3] where b\* is the modified branching factor given by the heuristic, which must be experimentally determined, and can approach 1 for very accurate heuristics, we can do a similar experiment to determine time complexity on our grid, O(d\*n), where d\* must be experimentally determined. In addition to the number of nodes visited, we must also consider the cost of adding and removing nodes from the heap, which in our case is a binary heap, and so has O(log(m)) time complexity, where m is the number of nodes added, d\*n. This means our total time complexity is O(d\*nlog(d\*n)). We choose to statically allocate a closed set equal to the size of the graph, so our memory complexity is O(V), where V is the number of vertices, if this became a limiting factor we could instead use a hashmap, reducing our memory complexity to O(d\*nlog(d\*n)), the same as our time complexity.

Both of our optimizations reduce the number of nodes that must be searched by a constant factor, this does not change our overall time complexity, so we expect a constant increase in speed from each optimization, rather than an increase that grows with the size b\*.

Lookahead uses strictly less memory than A\*. The memory usage of pruning depends on whether we’re permitted to modify the graph in-place or must allocate a new graph, if we must allocate a new graph then the memory complexity of pruning is O(V), where V >> d\*nlog(d\*n).

Pruning has an additional time complexity factor: the pruning process itself, which is O(V), which will be much larger than our pathfinding complexity. In a real world situation we would likely do many paths, which would amortize this one time cost. However if we truly only wanted one path, pruning could potentially cost us much more time than it saves. We will compare both the cost of pruning + pruned A\*, and the cost of pruned A\* alone, to measure these two cases. A secondary factor in pruning is the fact that it is a very simple process, this means that it could be offloaded to the gpu in a real world scenario, which could trivialize the cost.

**Graph setup**

For simplicity, we choose a grid of dimension 2, and we choose to path from the top left corner of the grid to the bottom right corner, this make the solutions easy to visualize and allows us to profile very large grids without running out of memory. Generalizing to higher dimensions and arbitrary dimensions should be trivial.

We chose to use Perlin noise to generate an initial distribution of walls on the graph. We then set the start and end tiles to be passable, and we use random sampling and flood fill to punch holes in the walls until all empty cells are connected.

**Results**

We compared the performance of A\*, lookahead, pruning, and lookahead with pruning on various grid sizes. We also recorded the number of nodes visited for each algorithm, to determine whether our assumptions about lookahead and pruning being constant factor reductions in number of nodes visited is true.

**Discussion**

**Further optimizations**

We talked about reducing memory usage by replacing the closed set array with a hashmap, we could also consider further reduce the number of heap pushes in lookahead by replacing the next variable with a stack, this would allow us to process all nodes whose heuristics are smaller rather than just one, resulting in less nodes pushed to the heap. A final optimization we could make would be to replace the heap with a list of stacks, a stack at index n would store nodes with f = n, this reduces the (amortized) time complexity of pushing and popping from O(log(n)) to O(1).

**References**

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